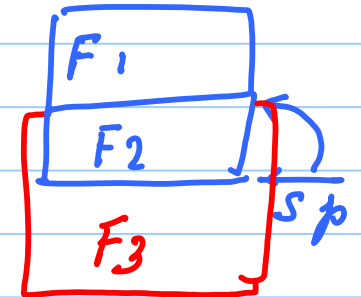
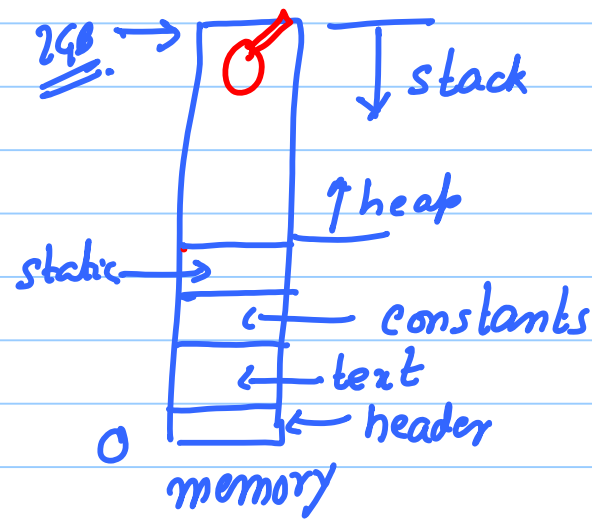


# Aug - 8

Note Title

08-08-2012

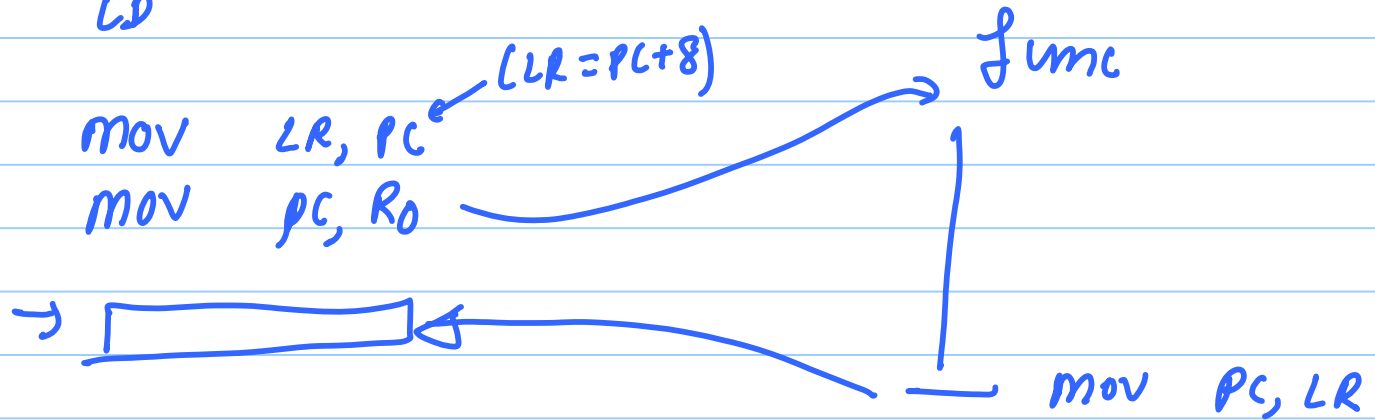
$$\left. \begin{array}{l} \textcircled{I} \text{ oldPC} = \text{PC} + 8 + \text{offset} \\ \textcircled{\text{book}} \text{ oldPC} = \text{PC} + 8 + \text{offset} \times 4 \end{array} \right\}$$



# function pointer

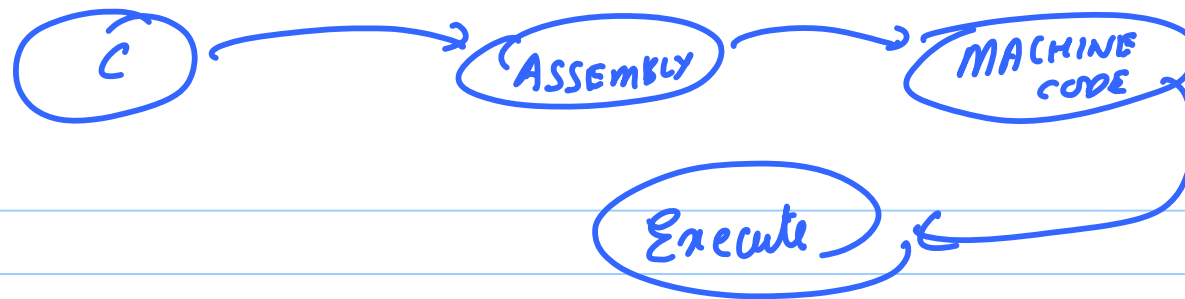
mov [R0] (...)  
LD

mov LR, PC (LR = PC + 8)  
mov PC, R0



ADD R3, R1, R2

MUL R3, R1, R2



- ✓ 1. (Adder/multiplier/divider)
- ✗ 2. Shifter/Logic
- ✓ 3. Memory

# Adders

$$a + b \rightarrow S$$

$$\text{Sum} \rightarrow a \oplus b$$

$$\text{Carry} \rightarrow a \cdot b$$

Half adder

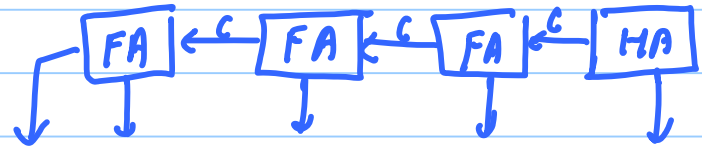
$$a + b + c \xrightarrow{\text{Carry}} \text{Sum}$$

$$\text{Sum} \rightarrow a \oplus b \oplus c$$

$$\text{Carry} \rightarrow ab + bc + ca$$

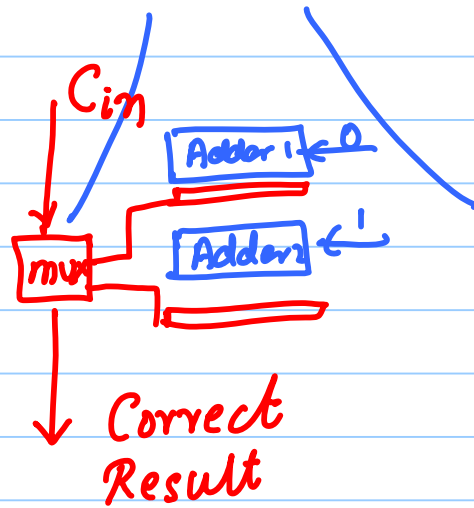
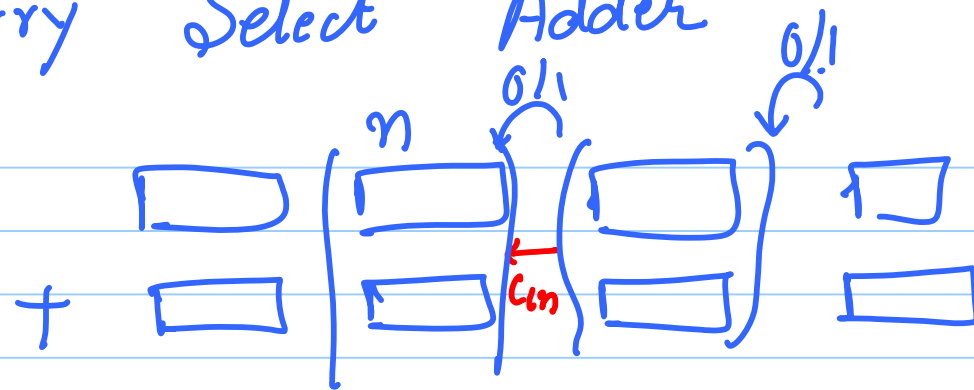
Full Adder

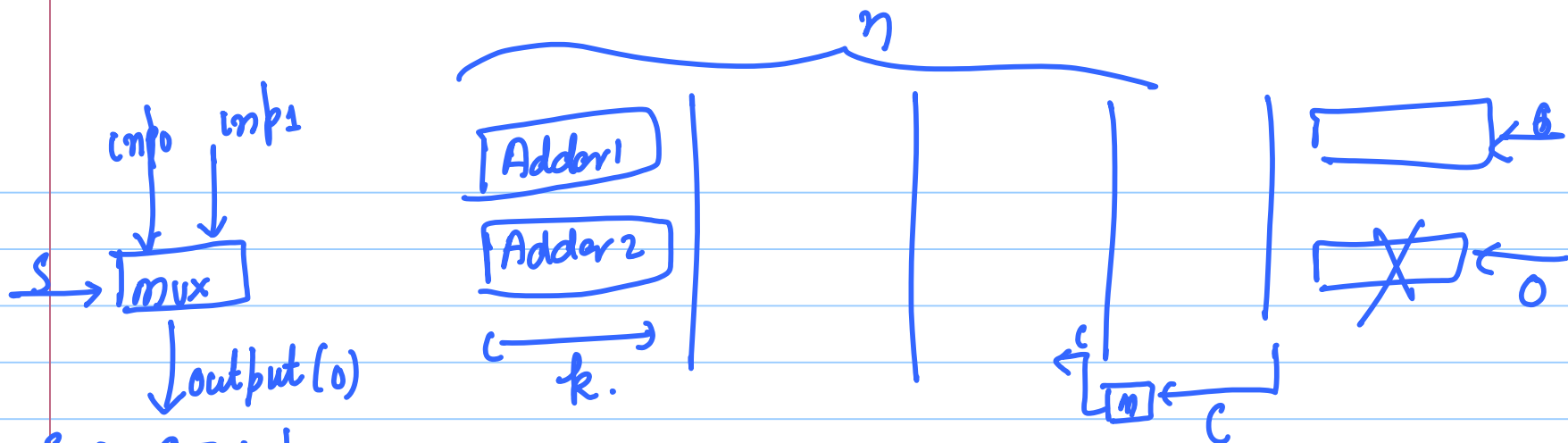
$$\begin{array}{r} \phantom{+} 1011 \\ + 1110 \\ \hline 11001 \end{array}$$



Ripple Carry Adder:  $O(n)$

# Carry Select Adder





$S=0, O = inp_0$   
 $S=1, O = inp_1$

Time per chunk :  $k$

Round 1:

Add each chunk. Produce two results

$C_{out_0} \leftarrow res_0$   
 $\quad \quad \quad 2$   
 $C_{out_1} \leftarrow res_1$

Round 2: Generate  $C_{in}$  for each chunk:  $(n/k)$

$$\text{Total Time: } k + n/k = O(\sqrt{n})$$

$$\frac{\partial (k + n/k)}{\partial k} \Rightarrow 1 - n/k^2 = 0$$

$$\Rightarrow n = k^2$$

$$\Rightarrow k = \sqrt{n}$$

Round 2:

$C_{in}^i$  (input carry of chunk  $i$ )

$C_{out}^j$  (output carry of chunk  $j$ )

$C_{out0}^i$  (output carry of adder 0 - chunk  $i$ )

$C_{out1}^i$  (output carry of adder 1  $\rightarrow$  chunk  $i$ )

$$C_{in}^{i+1} = f(C_{in}^i, C_{out0}^i, C_{out1}^i)$$

$$= C_{out1}^i C_{in}^i + C_{out0}^i \overline{C_{in}^i}$$

## Tighter Carry-Select Adder

$m$

$2$   $1$

$$\sum_{i=1}^m i = n$$

$$\Rightarrow \frac{m(m+1)}{2} = n.$$

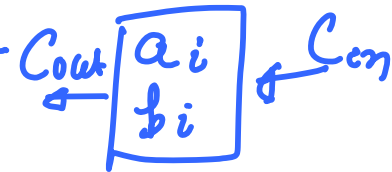
$$\Rightarrow m^2 - 2n + m = 0$$

$$m = \frac{\sqrt{1+8n} - 1}{2} = -\frac{1}{2} + \sqrt{\frac{1}{4} + 2n}$$

$$\approx \sqrt{2} \cdot \sqrt{n}$$



# Carry Lookahead Adder



Propagate (P)  $\rightarrow C_{out} = C_{in}$

$$\text{if: } P = a_i \oplus b_i = 1$$

Generate (G)  $\rightarrow a_i \cdot b_i$

$$C_{out} = G + P \cdot C_{in}$$



$$\begin{array}{r|l} a_2 & a_1 \\ + & b_2 & b_1 \\ \hline P_2 G_2 & P_1 G_1 \end{array}$$

$$G_{12} = G_2 + P_2 G_1$$

$$P_{12} = P_1 P_2$$

3-bit

$$P_{13} = P_1 P_2 P_3$$

$$G_{13} = G_3 + P_3 G_2 + P_3 P_2 G_1$$