CSL 630, Tutorial Sheet 5

- 1. Give an O(n+m) time i.e. linear time algorithm to determine if a 2-SAT formula is satisfiable. The formula has n variables and m clauses. (Hint: $(\mathbf{x} \mathbf{V} \mathbf{y})$ is equivalent to $(\bar{x} \to y), (\bar{y} \to x)$.)
- 2. Prove that the following problems are NP-complete:

(a) **CLIQUE**

Instance: Graphs G=(V,E) and an integer $k \leq |V|$.

Question: Is there a clique of size at least k in G?

(a) SUBGRAPH ISOMORPHISM

Instance: Graphs G=(V,E) and H=(V',E').

Question: Is H isomorphic to a subgraph of G?

Definition: Graph $G_1 = (V_1, E_1)$ is **isomorphic** to $G_2 = (V_2, E_2)$ if there is a one-one onto function $f: V_1 \to V_2$ such that $(u, v) \in E_1$ iff $(f(u), f(v)) \in E_2$.

(b) **DOMINATING SET**

Instance: Graph G=(V,E) and positive integer k.

Question: Is there $V' \subseteq V, |V'| = k$ such that each vertex u in (V-V') is adjacent to some vertex v in V'.

- 3. (a) Show that **PARTITION** is self-reducible, i.e. give a polynomial time algorithm to solve the search problem , given a subroutine for the decision problem.
 (b) Show that VERTEX COVER and CLIQUE are self reducible.
- 4. Show the existence of a co-NP complete problem. Provide all formal definitions.