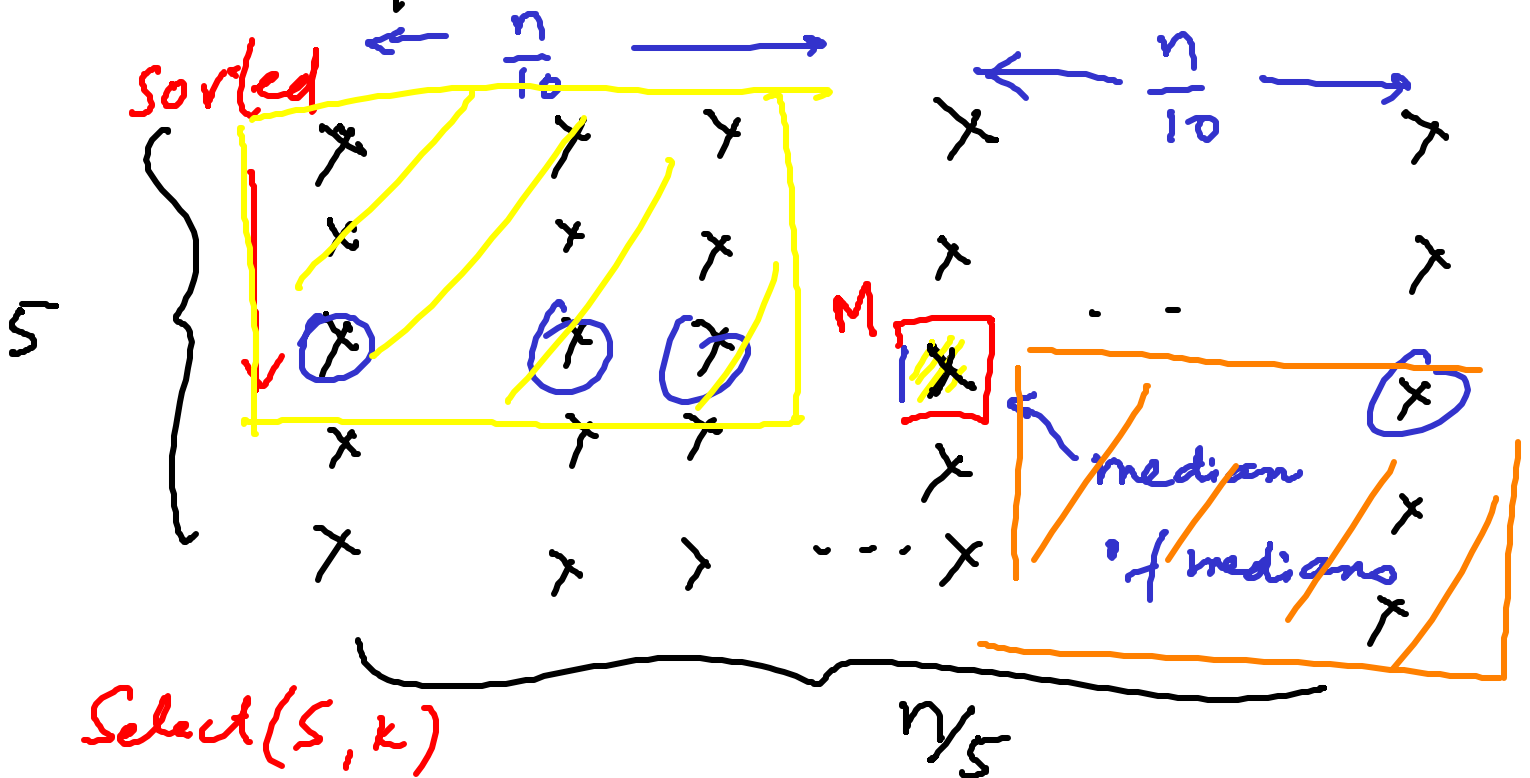


Deterministic Selection

How to find a "good" splitter,
 so that at least a constant fraction,
 say $\frac{n}{4}$ elements are eliminated



Select(S, k)

$\frac{n}{5}$

- Sort every column containing at most 5 elements

$$\text{Time} : O(1) \cdot \frac{n}{5} = O(n)$$

- Find the "median" of the median of each row, i.e. median of $\frac{n}{5}$ elements
- Find the rank of M $O(n)$

Claim: The median of medians
has rank in the range $[\alpha n, (1-\alpha)n]$
for some constant $\alpha < 1$

$\frac{n}{10} \times 3$ elements are smaller than M

$\frac{3n}{10}$ " " larger than M

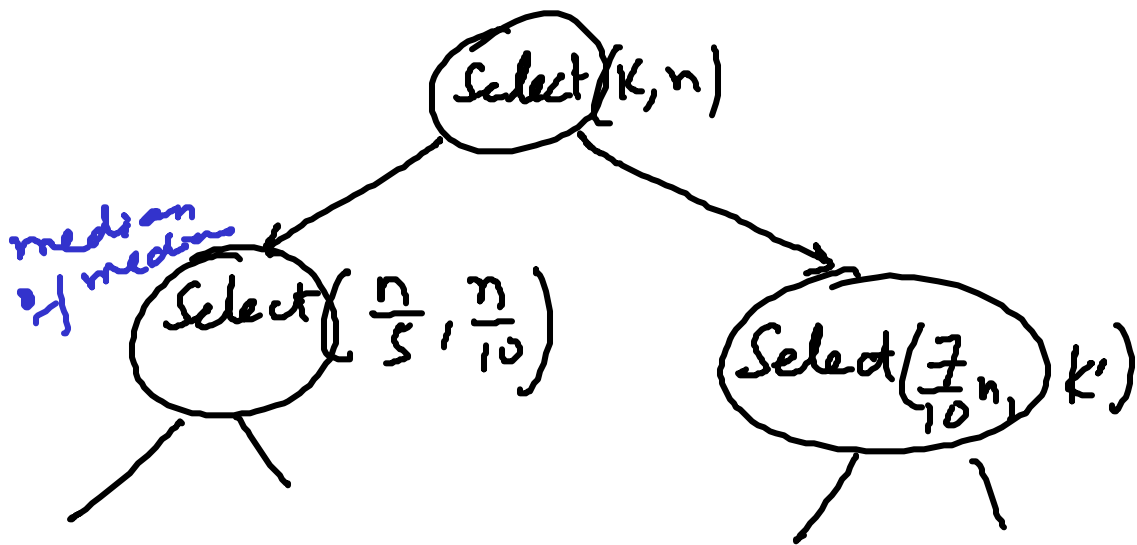
$$\alpha = \frac{3}{10}$$

Step 4: If $\text{rank}(M, S) < k$
then we can discard all element
in the shaded box to the right bottom

Time: $O(n)$

Step 5: Call recursively on the
remaining elements with the
adjusted value of k .

$T(n)$: the worst case time over all $k \leq n$
 $T(n) = T\left(\frac{n}{5}\right) + O(n) + T\left(\frac{7n}{10}\right)$ $T(2) = O(1)$



c. Prove (using induction) that

$$T(n) \leq \alpha \cdot n \text{ for some constant } \alpha.$$

Soft-heaps by Chazelle

Prune and Search

Finite alphabet Σ

Strings s_1, s_2, \dots, s_n defined
over Σ of lengths l_1, l_2, \dots

$$|s_i| = l_i$$

01, 1100, 101010, ...
...

$$L = \sum l_i$$

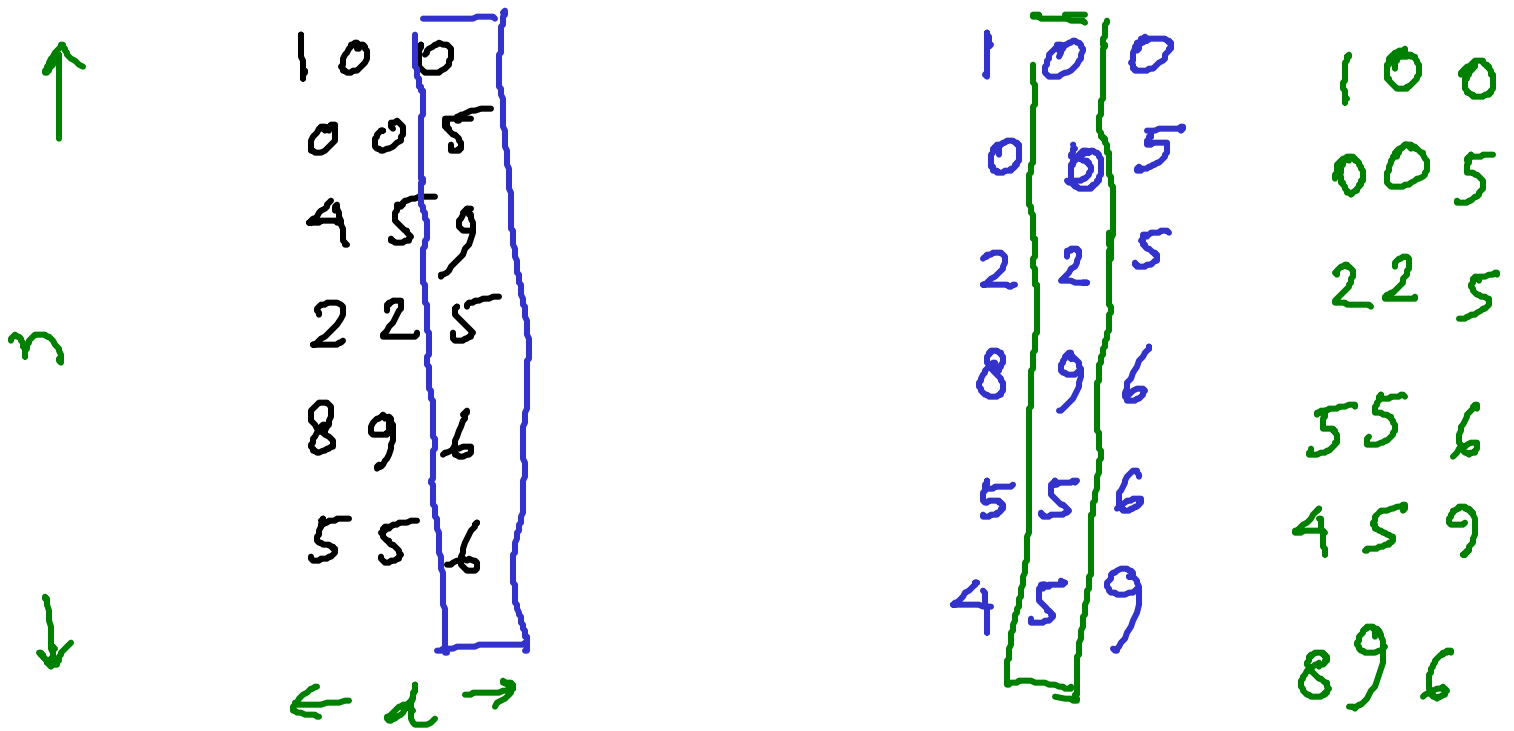
possible that
some string has size 1
whereas some string
may have size $\frac{L}{4}$

We would like to order the strings
use "lexicographic ordering"

What algorithm should we use?

Radix sort?

100, 255, 099, 005, ... - 855



0 0 5
1 0 0
2 2 5
4 5 9
5 5 6
8 9 6

Using count sort,
time for each
phase
 $\Theta(|\Sigma| + n)$
range

Stable sort: doesn't alter the ordering of numbers with identical

Radix sort has to be based on some stable sorting

Count Sort : Stable?

Every comparison sorting can be made stable

Radix Sort \checkmark^n

$$O(d (\sum 1 + n))$$

As long as d is constant

$$\sum = 1..n, \text{ i.e. } \log n \text{ bits}$$

We can sort n numbers in
the range $[1..n^d]$ in

$$O(n) \text{ time.}$$