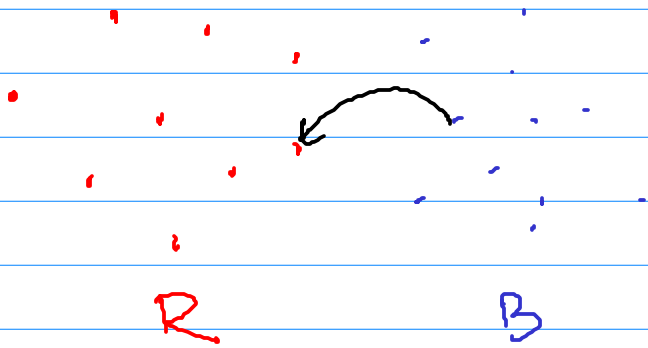


# Lecture 40: Shape Comparison



$$|R| = |B| = n$$

$$\rightarrow Q: B \rightarrow R$$

Directed Hausdorff distance

$$H(B, R) = \max_{b \in B} \|b - Q(b)\|$$

:  $L_\infty$ -norm distance

$M(R, B)$

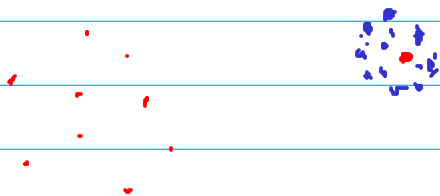
$$\sum_{b \in B} \|b - Q(b)\|$$

:  $L_1$ -norm

$$\sqrt{\sum_{b \in B} \|b - Q(b)\|^2}$$

$L_2$ -norm

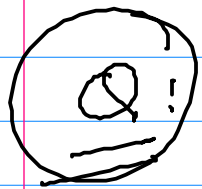
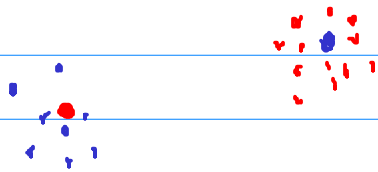
✓  $Q(b)$ : Nearest neighbor of  $b$  in  $R$   
 $O(n \log n)$  time in  $\mathbb{R}^2$



$$H(B, R) = \max_{b \in B} \min_{r \in R} \|b - r\|$$

Hausdorff distance  $\underbrace{\min_{r \in R} \|b - r\|}_{\|b - Q(b)\|}$

$$h(R, B) = \max \{H(B, R), H(R, B)\}$$



Q! bijective

1-1 mapping between R & B

Minimum weight matching between R & B.

$$G = (V, E): w: E \rightarrow \mathbb{R}^+$$

Matching:  $M \subseteq E$ : vertex disjoint edges

$G = (R \cup B, R \times B)$  bipartite graph

$$w(r, b) = \|r - b\|$$

$$O(mn)$$

$$O(n^3) \text{ algorithm } m \geq n^2$$

For  $(RUB, R \times B)$ :

best-known algorithm takes

$\tilde{O}(n^2)$  time

Is there a subquadratic algorithm?

$\sim n^{1+1/c}$  time

logc approximation!

$O(1)$  factor - estimate the cost  
in  $O(n \log^2 n)$

Earth Mover's distance (transportation distance)

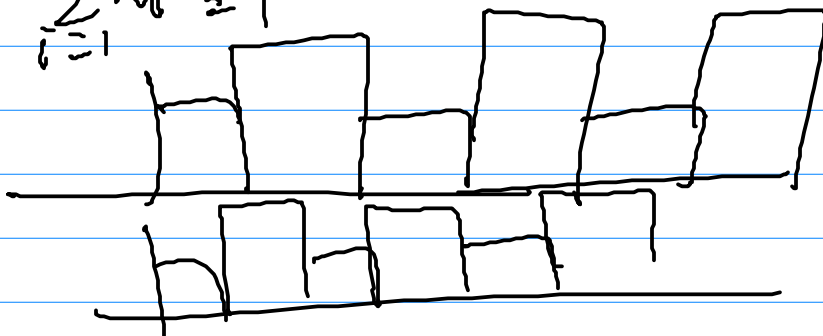
$R, B = \{b_1, \dots, b_n\}$

$\{r_1, \dots, r_n\}$   $w_1' \dots w_n'$

$w_1 \dots w_n$

$$\sum_{i=1}^n w_i = 1$$

$$\sum_{i=1}^n w_i = 1$$



$R, B$  and given  $\mu$

Find  $\mathbb{T}$ : set of all rigid motions

compute

$$\min_{\tau \in \mathbb{T}} \mu(R, \tau(B))$$

ICP: Iterative Closest Pair Algorithm

$B_i$ : current position of  $B$

$R$

compute

$Q(b)$  for all  $b \in B_i$

is fixed



Translate  $B_i$  s.t.

$$\tau^* = \arg \min_{\tau \in \mathbb{T}} \sum_{b \in B} \|b - Q(b) \circ \tau\|^2$$

$$B_{i+1} = \tau^*(B_i)$$