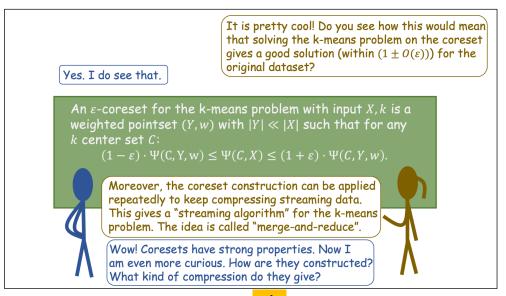
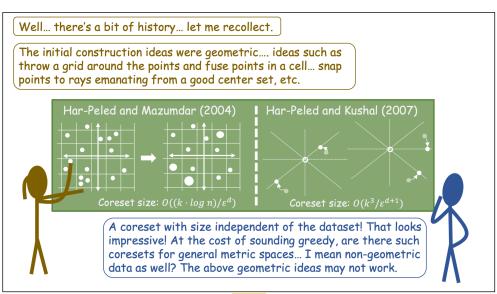


Okay... let me define an ε -coreset for the k-means problem.

An ε -coreset for the k-means problem with input X,k is a weighted pointset (Y,w) with $|Y|\ll |X|$ such that for any k center set C: $(1-\varepsilon)\cdot \Psi(C,Y,w)\leq \Psi(C,X)\leq (1+\varepsilon)\cdot \Psi(C,Y,w).$ I see... so a coreset (Y,w) mimics the dataset X in terms of retaining the k-means cost for "any" k-center set. That's awesome!



3



For general metric spaces, there are sampling-based constructions. They use ideas such as sample from concentric "rings" around a representative k center set and reweight the sampled points to represent points in the same ring.

Unlike the geometric coreset, I see a dependence on the data size n.

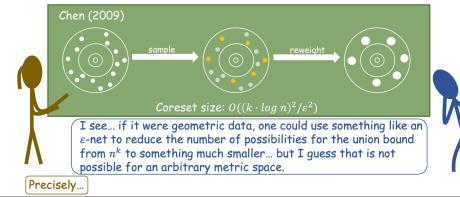
Chen (2009)

Coreset size: $O((k \cdot log \, n)^2/\epsilon^2)$ Oh ...that comes from the union-bound...

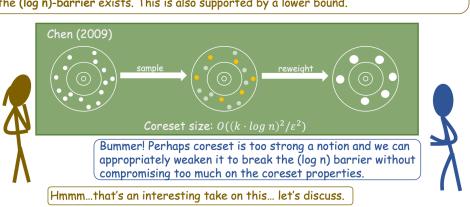
let me explain.

The Metric k-means problem: Let (X, D) be a metric. Given $X \subset X$ and k, find a center set $C \subset X$ with |C| = k such that the cost $\Psi(C, X) \equiv \sum_{x \in X} \min_{n \in C} D(x, c)^2$ is minimized.

It's a randomized construction. The argument works like this: for a fixed k center set C, (sampling+reweighting) $(k/\varepsilon)^2$ points satisfies the coreset property with high probability. However, we want it to work for "every" k center set and there could be n^k of them. For the union-bound to work, we must blow up the sample size by a factor of (log n).

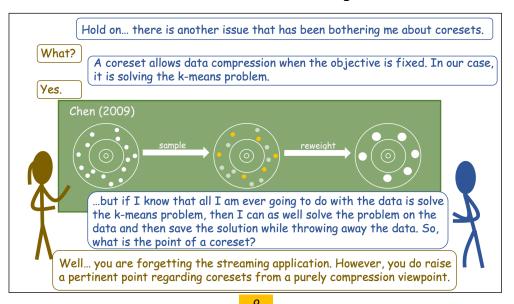


In fact, for specific metric spaces, such as the Euclidean space, these ε -set, ε -centroid set ideas have been used to obtain small sized coresets... and that has been the agenda of most research on coresets in the recent past. However, for general metric spaces, the (log n)-barrier exists. This is also supported by a lower bound.



7

3



Discussions on coresets have lately drifted completely towards the streaming application, even though it was initially designed as a data compression tool.

I have an application in mind that may justify the data compression aspect. Lately, constrained clustering has gained importance where the clustering must satisfy certain constraints in addition to optimizing the k-means cost.

Constrained clustering:

Balanced k-means clustering

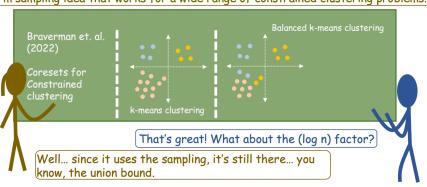
One example is balanced clustering, where the clusters should have roughly equal size. Other examples include various fairness notions.

I'm aware... could you elaborate on the coreset aspect?

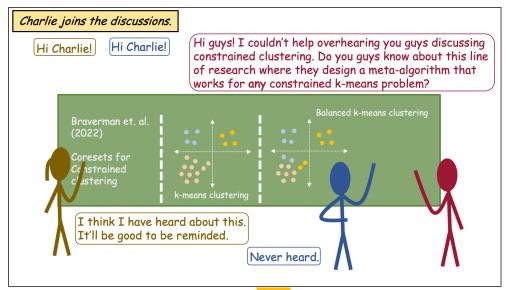
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If we can compress while being oblivious of the specific constraints the clusters need to satisfy, then computing the coreset makes sense. Many times, the constraints are not known at the time of access to the data.

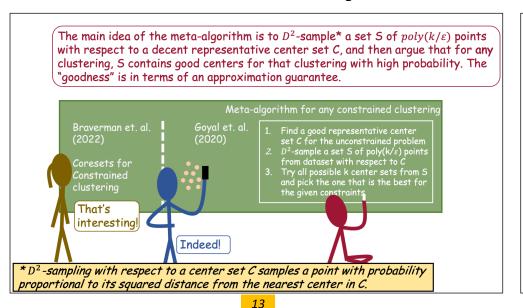
You are in luck! There has been a recent development in coreset construction using Chen's uniform sampling idea that works for a wide range of constrained clustering problems.

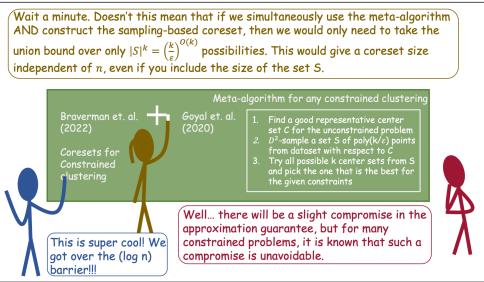


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12





14

This new coreset notion, even though interesting, raises many more interesting questions, including one about streaming construction. That's true. However, I'm excited! What should we call our coreset? How about "Universal Weak Coreset?" Meta-algorithm for any constrained clustering Braverman et. al. 1. Find a good representative center That's interesting. "Universal" indicating the fact Coresets for that it works for many constrained problems, and I Constrained guess "Weak" because it reminds one of the weak coreset idea from the Euclidean setting. Yes. Should we publish? Well... publication is always tricky ... ours is a conceptual contribution ... people may want to see technical development ... anyway, I think our new notion may advance further research on coresets... so we should publish.

