

Consider the following 3-partition problem. Given integers a_1, a_2, \dots, a_n , we want to determine whether it is possible to partition $\{1, 2, \dots, n\}$ into three disjoint subsets I, J, K such that

$$\sum_{i \in I} a_i = \sum_{j \in J} a_j = \sum_{k \in K} a_k = \frac{1}{3} \sum_{i=1}^n a_i$$

For example, for input $(1, 2, 3, 4, 4, 5, 8)$ the answer is *yes* because there is the partition $(1, 8), (4, 5), (2, 3, 4)$. On the other hand, for input $(2, 2, 3, 5)$ the answer is *no*. Give an efficient algorithm for this problem. (10)

We solve this by Dynamic Programming. Let X be a 3-dimensional array of size $n \times W \times W$ where $W = \frac{1}{3} \sum_{i=1}^n a_i$.

Define $X[i, j, k] = 1$ if $\exists S_1, S_2 \subseteq \{a_1, a_2, \dots, a_i\}$, $S_1 \cap S_2 = \emptyset$
and $\sum_{a_i \in S_1} a_i = j$ & $\sum_{a_i \in S_2} a_i = k$.
 $= 0$ otherwise.

Now,

$$X[i, j, k] = 1 \text{ if } X[i-1, j, k] = 1 \text{ OR } X[i-1, j-a_i, k] = 1 \\ \text{OR } X[i-1, j, k-a_i] = 1.$$

Hence the entire array can be filled in time $O(nW^2)$.

Finally, the answer to the 3-partition problem is YES if $X[n, \frac{W}{3}, \frac{W}{3}] = 1$

COMMON MISTAKES: Many students have given the following solution:
Find a subset S whose sum is W , remove these integers & find another subset of sum W .

For the instance $(1, 2, 3, 4, 4, 5, 8)$ if the first subset found is $(1, 3, 5)$ then the remaining integers are $(2, 4, 4, 8)$ & there is no subset here of sum 9 . However the original instance is a YES instance as shown above (in the question).

There is no way of formulating this as a max flow problem. Some students have created networks with edge capacity equal to the given integers. There is no way to ensure that these edges would be saturated.

MARKING SCHEME:

2 marks given if you gave the "wrong solution" described above
i.e. find a set S of sum W , remove it & repeat.