

**Important:** The question marked with a ♠ is to be submitted via gradescope by 11:59PM on the day that you have your tutorial.

**Problem 1 [1, Prob 12, page 30]**

Show that every 2-connected graph contains a cycle.

**Problem 2**

Show using only the material covered in [1, Ch 1.4] that every connected graph on  $n$  vertices has at least  $n - 1$  edges.

**Problem 3**

Generalize the result of Problem 2 to show that every graph on  $n$  vertices and  $m$  edges has at least  $n - m$  components.

**Problem 4 ♠**

Given a graph  $G = (V, E)$  and a minimal cut  $F \subseteq E$ , show that any cycle of  $G$  contains an even number of edges of  $F$  (this number could be 0 as well).

**Problem 5**

Show that if there is a vertex  $v$  of odd degree in graph  $G$  there must be a path from  $v$  to another vertex  $u$  of  $G$  which also has odd degree.

**Problem 6**

Let  $\bar{G}$  be the complement of the graph  $G$ , i.e., all edges of  $G$  are non-edges of  $\bar{G}$  and vice versa. Show that both  $G$  and  $\bar{G}$  cannot be disconnected, i.e., at least one of them must be connected.

**Problem 7**

Given a graph  $G = (V, E)$  such that  $|V| = n$ , a cut  $F \subset E$  is called a *balanced cut* if  $G \setminus F$  has exactly 2 components and each of these components has size at least  $n/3$ . Construct graphs on  $n$  vertices whose smallest balanced cut has size (a)  $\theta(1)$ , (b)  $\theta(\sqrt{n})$  and (c)  $\theta(n)$ .

**Problem 8 (Menger's Theorem)**

Prove that a graph  $G$  has  $\lambda(G) = k$  for any  $k \geq 1$  iff there are  $k$  edge-disjoint paths between any pair of vertices in  $G$ . Two paths are said to be independent if they don't share any edges. Caution: One direction of this theorem is easy and the other is tricky.

## References

- [1] Reinhard Diestel, *Graph Theory 5ed.*, Springer, 2016.